



TECHNISCHE  
UNIVERSITÄT  
DRESDEN



HAEC

# SEMANTIC TECHNOLOGY FOR CONTEXT AWARENESS

CRC “Highly Adaptive Energy-efficient computing” (HAEC)

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# Motivation

## Employing ontology-based situation recognition

- support context awareness by logical reasoning  
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- enrich raw data by semantic technologies
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- detect high-level, composite events

# Our Setting

in the HAEC project

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We: recognize situations (when adaptations might be beneficial)

Others: perform adaptation

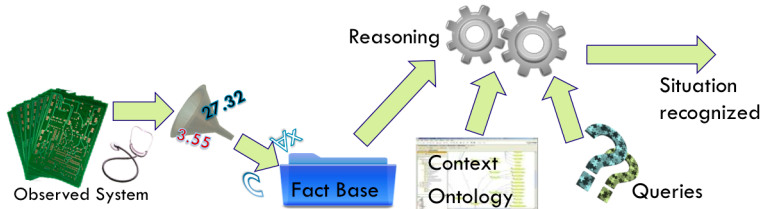
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### The Setting

- **Background ontology:** description logics (DL) TBox
- **Input data:**
  - information on components of observed system (static), sensor data (dynamic)
  - stored in DB, connected to TBox via preprocessors (/mappings)
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Essentially an OBDA setting, but

- **temporal sequence** of system states: data items time stamp annotated.
- data items grouped into logical categories with a **membership degree**

# Temporal Query Answering

## Motivation and Approach

- System is observed over time yielding sequence of system snapshots.
- “growing” window:  
storing (relevant) data from start until current time point
- Temporal query language: extension of DL by LTL to navigate on the temporal sequence

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### Temporal conjunctive queries (TCQs)

- Combine LTL and DLs
- built inductively from CQs using:  $\neg, \wedge, \bigcirc^-, \diamond^-, \square^-$

# Temporal Query answering

## Results

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- 2-step rewriting: uses atemporal rewriting as black box procedure

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Horn-DLs (include DL-Lite,  $\mathcal{EL}$ )

- with fixed set of queries ( $\sim$  temporal DB monitoring problem): bounded history encodings
- treatment of future operators: gain of one exponent

# Fuzzy Queries

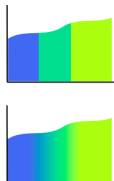
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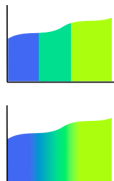




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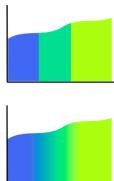
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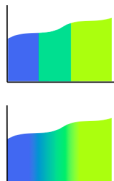


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Queries with membership degrees:

- fuzzy query: degrees per atom
- threshold query: degree per CQ

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- data complexity as in the crisp case!

# Future Work

## What's next?

Theoretical side:

- incorporate **probabilities** for the data items  
e.g., cope with erroneous sensors

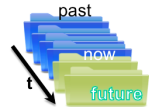
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Practical side:

- exploit **parallel hardware**
- investigate **incremental reasoning**